School of Computer Science & Engineering

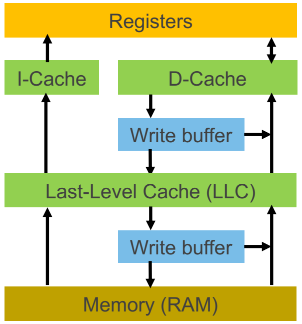
COMP9242 Advanced Operating Systems

2020 T2 Week 03b

Caches:

What Every OS Designer Must Know

@GernotHeiser



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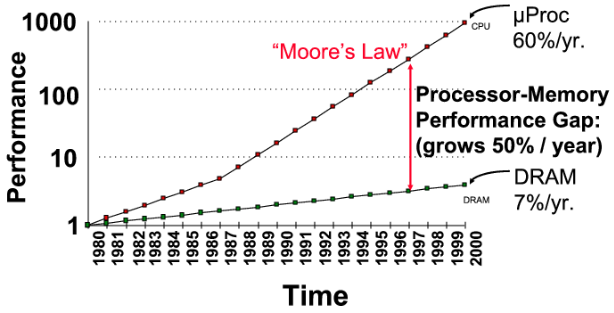
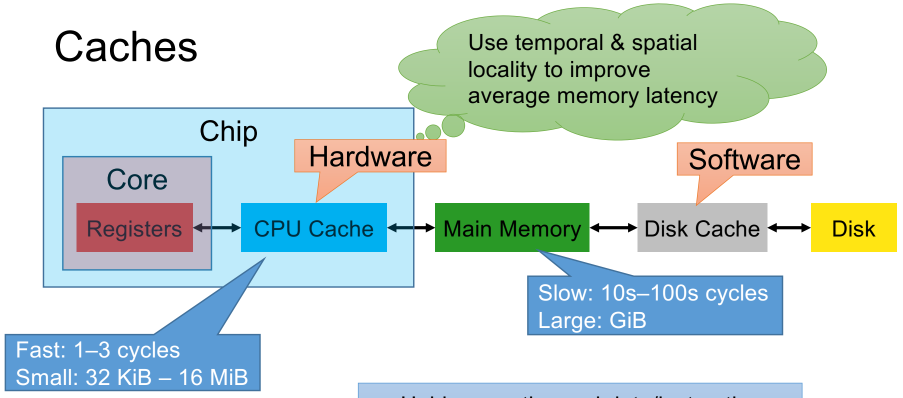
The Memory Wall



Multicore offsets stagnant per-core performance with proliferation of cores

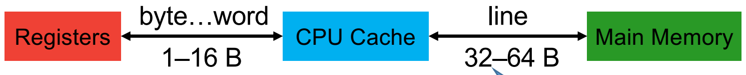
- Same effect on overall memory bandwidth
- Basic trend is unchanged

Caches



- Holds recently used data/instructions
- Load/fetch hits in cache ⇒ fast access
- Miss not much worse than no cache
- Key is high hit rate (>90%)

Cache Organisation: Unit of Data Transfer



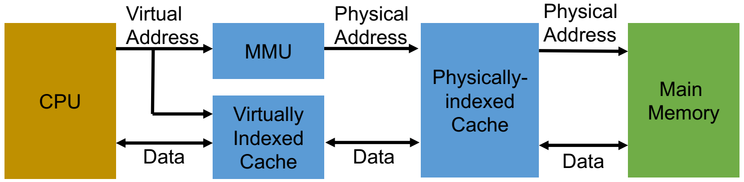
Line is also unit of allocation, holds data and

- valid bit
- modified (dirty) bit
- tag
- access stats (for replacement)

Reduce memory transactions:

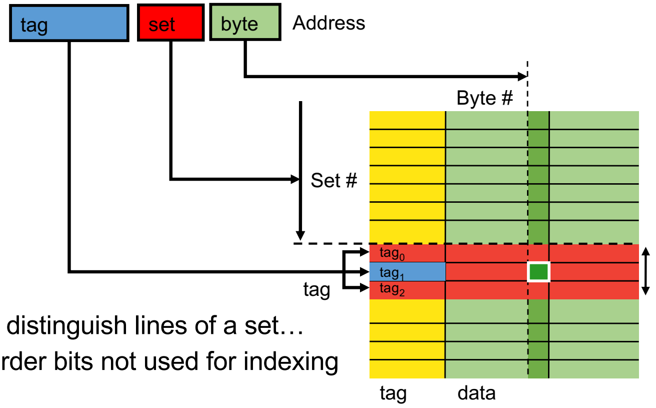
- Reads – locality
- Writes – clustering

Cache Access



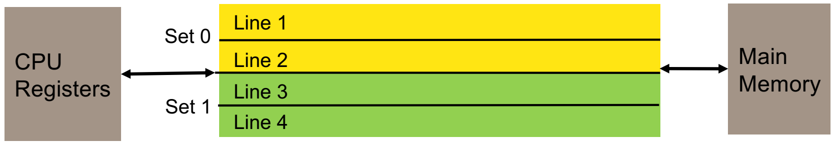
- Virtually indexed: looked up by virtual address
  - operates concurrently with address translation
- Physically indexed: looked up by physical address
  - requires result of address translation
- Usually a hierarchy: L1, L2, …, LLC (last-level cache, next to RAM)
  - L1 may use virtual address, all others use physical only

Cache Indexing



- The *tag* is used to distinguish lines of a set…
- Consists of high-order bits not used for indexing

Cache Indexing

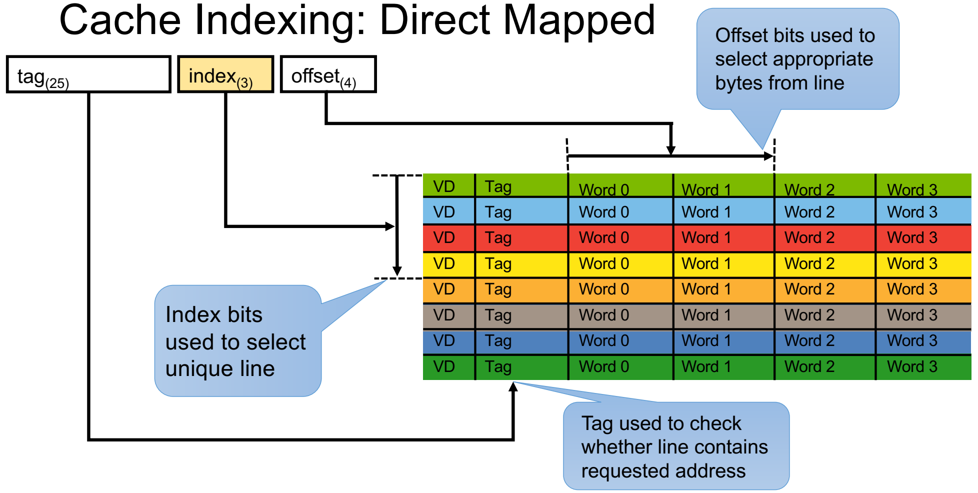


- Address hashed to produce index of line set
- Associative lookup of line within set
- n lines per set: n-way set-associative cache, typically n=1–16
  - n = 1 is called direct mapped
  - 2 ≤ n ≤ ∞ is called *set associative*
  - n = ∞ is called fully associative
- Hashing must be simple (complex hardware is slow)
  - generally use least-significant bits of address

Many conflicts ⇒ low hit rate

Slow & power-hungry

Cache Indexing: Direct Mapped



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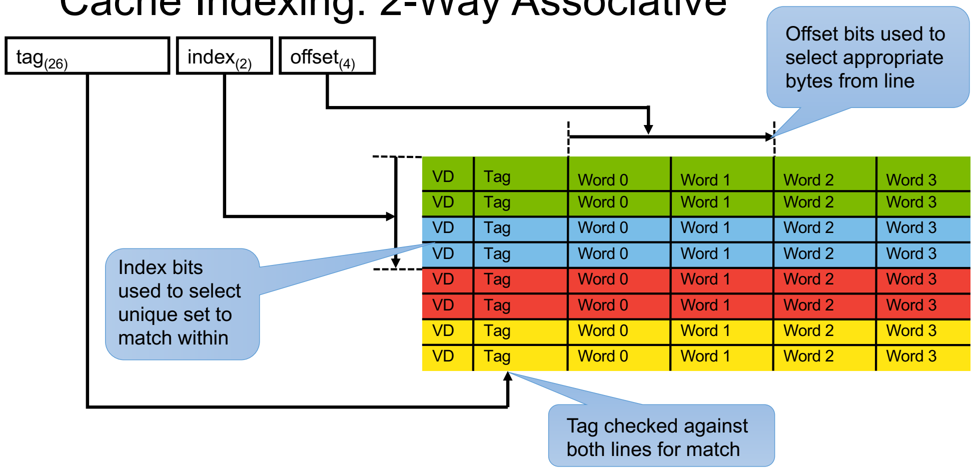
Offset bits used to select appropriate bytes from line

Index bits used to select unique line

Tag used to check whether line contains requested address

Cache Indexing: 2-Way Associative



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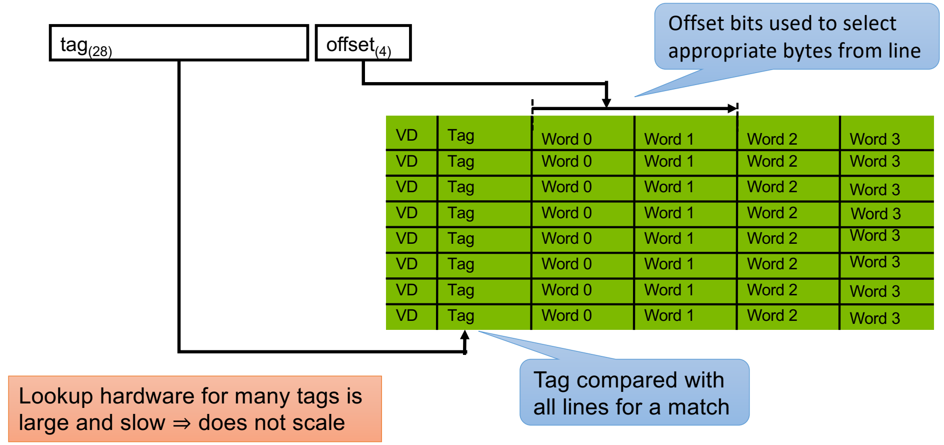
Offset bits used to select appropriate bytes from line

Index bits used to select unique set to match within

Tag checked against both lines for match

Cache Indexing: Fully Associative



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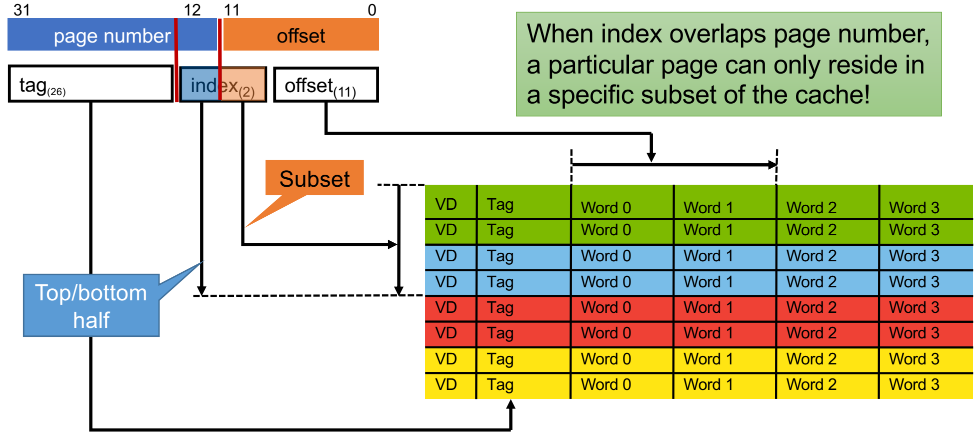
Offset bits used to select appropriate bytes from line

Tag compared with all lines for a match

Lookup hardware for many tags is large and slow ⇒ does not scale

Cache Associativity vs Paging



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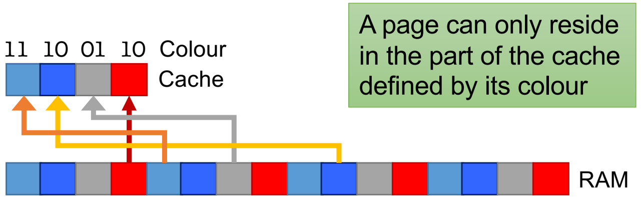
When index overlaps page number, a particular page can only reside in a specific subset of the cache!

Cache Mapping Implications

Multiple memory locations map to same cache line



A page can only reside in the part of the cache defined by its colour

If c index bits overlap page #, a page can only reside in 2^{-c} of the cache

Cache is said to have 2^c colours
2^c = cache_size/(page_size × assoc)

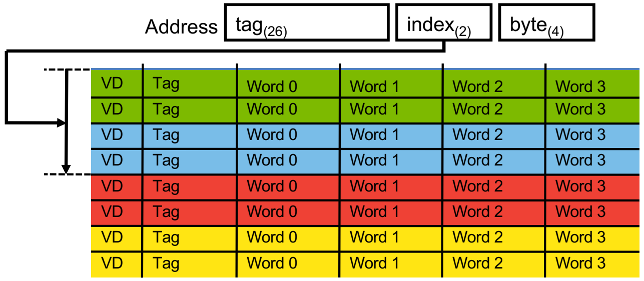
Cache Misses

- *n-way* associative cache can hold *n lines* with the same *index* value
- More than *n lines* are competing for same index forces a miss!
- There are different types of cache misses ("*the four Cs*"):
  - **Compulsory miss:** data cannot be in the cache (of infinite size)
    - First access (after loading data into memory or cache flush)
  - **Capacity miss:** all cache entries are in use by other data
    - Would not miss on infinite-size cache
  - **Conflict miss:** all lines *with the same index value* are in use by other data
    - Would not miss on fully-associative cache
  - **Coherence miss:** miss forced by hardware coherence protocol
    - Covered later (multiprocessing lecture)

Cache Replacement Policy

- Indexing (using address) points to specific line set
- On miss (no match and all lines valid): *replace* existing line
  - Dirty-bit determines whether write-back needed
- Replacement strategy must be simple (hardware!)

Typical policies:
- LRU
- pseudo-LRU
- FIFO
- "random"
- toss clean



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Cache Write Policy

- Treatment of store operations
  - **write back:** Stores only update cache;
    memory is updated once dirty line is replaced (flushed)
    - ☑clusters writes
    - ⌘memory inconsistent with cache
    - ⌘multi-processor cache-coherency challenge
  - **write through:** stores update cache and memory immediately
    - ☑memory is always consistent with cache
    - ⌘increased memory/bus traffic
- On store to a line not presently in cache (write miss):
  - **write allocate:** allocate a cache line and store there
    - typically requires reading line into cache first!
  - **no allocate:** store directly to memory, bypassing the cache

Typical combinations:
- write-back & write allocate
- write-through & no-allocate

Cache Addressing Schemes

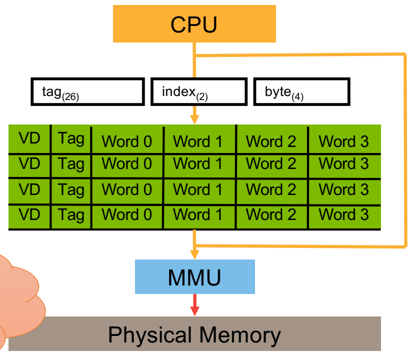
- So far pretended cache only sees one type of address: virtual or physical
- However, *indexing and tagging can use different addresses!*
- Four possible addressing schemes:
  - *virtually-indexed, virtually-tagged* (VV) cache
  - *virtually-indexed, physically-tagged* (VP) cache
  - *physically-indexed, virtually-tagged* (PV) cache
  - *physically-indexed, physically-tagged* (PP) cache

Nonsensical except with weird MMU designs

Virtually-Indexed, Virtually-Tagged Cache

- Also called *virtually-addressed cache*
- Various incorrect names in use:
  - ~~virtual cache~~
  - ~~virtual address cache~~
- Uses virtual addresses only
- Can operate concurrently with MMU
- Usable for on-core L1
  - Rarely used these days

Permissions?
Write back?



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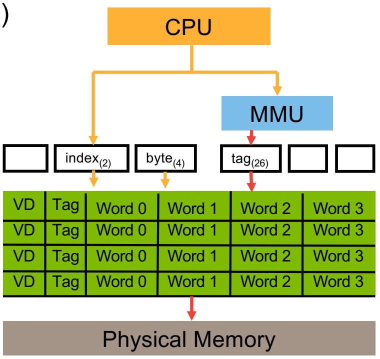
Virtually-Indexed, Physically-Tagged Cache

- Virtual address for accessing line (lookup)
- Physical address for tagging
- Needs complete address translation for looking up retrieving data
- Indexing concurrent with MMU
- Used for on-core L1

Use MMU for tag check & permissions



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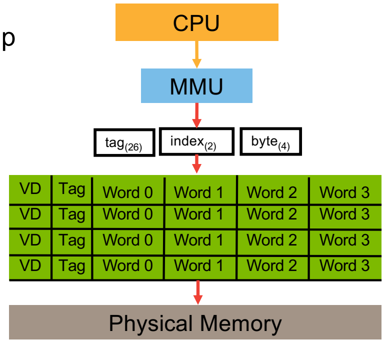
Physically-Indexed, Physically-Tagged Cache

- Only uses physical addresses
- Address translation result needed for lookup
- Only sensible choice for L2…LLC

Speed matters
less after L1 miss

Page offset invariant under VA→PA:
- Index bits ⊂ offset bits
  ⇒ don't need MMU for indexing!
- VP = PP in this case
  ⇒ fast, suitable for L1
- Single-colour cache!



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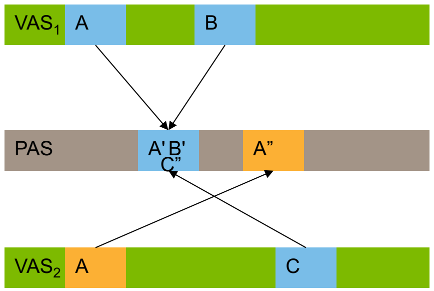
Cache Issues

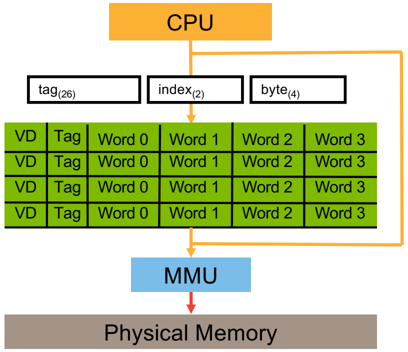
- Caches are managed by hardware transparently to software, so OS doesn't have to worry about them, ~~right?~~ **Wrong!**
- Software-visible cache effects:
  - *performance*
    - cache-friendly data layout
  - *homonyms*:
    - same address, different data
    - can affect correctness!
  - *synonyms (aliases)*:
    - different address, same data
    - can affect correctness!



Virtually-Indexed Cache Issues

Homonyms – same name for different data:

- Problem: VA used for indexing is context-dependent
  - same VA refers to different PAs
  - tag does not uniquely identify data!
  - wrong data may be accessed
  - an issue for most OSes
- Homonym prevention:
  - flush cache on each context switch
  - force non-overlapping address-space layout
    - single-address-space OS
  - *tag* VA with *address-space ID* (ASID)
    - makes VAs global



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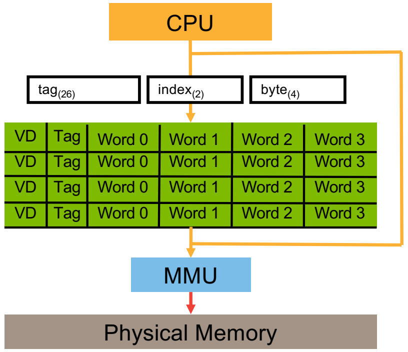
Virtually-Indexed Cache Issues

Synonyms – multiple names for same data:

- Several VAs map to the same PA
  - frame shared between ASs
  - frame multiply mapped within AS
- May access stale data!
  - same data cached in multiple lines
    - … if aliases differ in colour
  - on write, one synonym updated
  - read on other synonym returns old value
  - physical tags or ASIDs don't help!
- Are synonyms a problem?
  - depends on page and cache size (colours)
  - no problem for R/O data or I-caches



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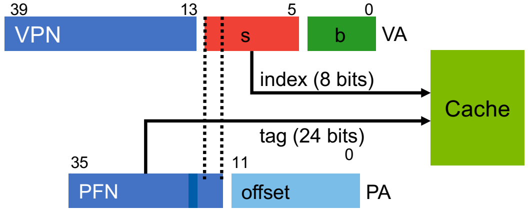
Example: MIPS R4x00 Synonyms

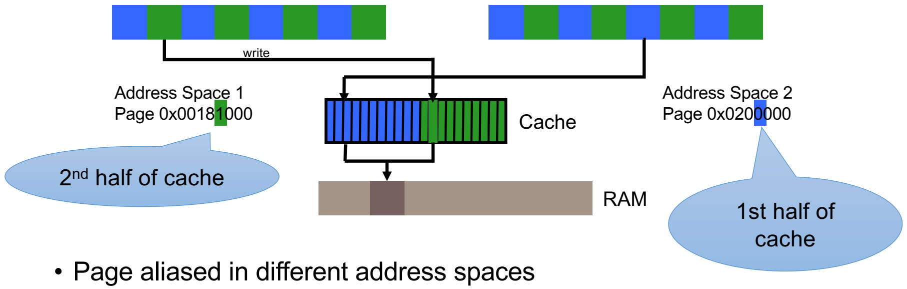
- ASID-tagged, on-chip VP cache
  - 16 KiB cache, 2-way set associative, 32 B line size, 4 KiB (base) page size
  - size/associativity = 16/2 KiB = 8 KiB > page size (2 page colours)
    - 16 KiB / (32 B/line) = 512 lines = 256 sets ⇒ 8 index bits (12..5)
    - overlap of tag bits and index bits, but from different addresses!

Remember, only index determines location of data!

- Tag only confirms hit
- Synonym problem iff $\text{VA}_{12} \neq \text{VA'}_{12}$
- Problem of virtually-indexed cache with multiple colours

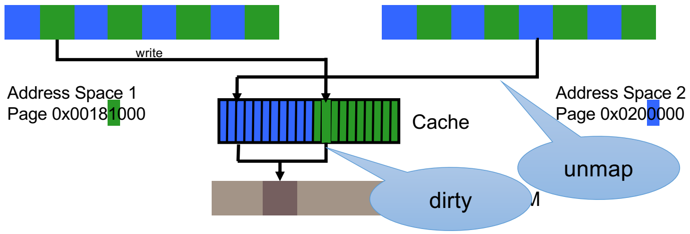


Address Mismatch Problem: Aliasing



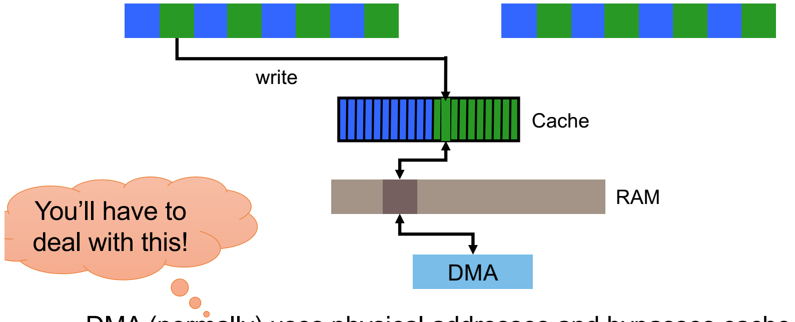
- Page aliased in different address spaces
  - AS_1: $VA_{12} = 1$, AS_2: $VA_{12} = 0$
- One alias gets modified
  - in a write-back cache, other alias sees stale data
  - *lost-update problem*

Address Mismatch Problem: Aliasing



- Unmap aliased page, remaining page has a dirty cache line
- Re-use (remap) frame for a different page (in same or different AS)
- Access new page
  - without replication, new write will overwrite old (hits same cache line)
  - with replication, alias may write back after remapping: "*cache bomb*"

DMA Consistency Problem



- DMA (normally) uses physical addresses and bypasses cache
  - CPU access inconsistent with device access
  - must flush cache before device write
  - must invalidate cache before device read

Avoiding Synonym Problems

- Flush cache on context switch
  - doesn't help for aliasing *within* address space!
- Detect synonyms and ensure:
  - all read-only, or
  - only one synonym mapped
- Restrict VM mapping so synonyms map to same cache set
  - eg on R4x00: ensure $VA_{12} = PA_{12}$ – colour memory!
- Hardware synonym detection
  - e.g. Cortex A53: store overlapping tag bits of both addresses & check
  - "physically"-addressed

Summary: VV Caches

- ☑Fastest (don’t rely on TLB for retrieving data)
  - ⌘still need TLB lookup for protection
  - ⌘… or alternative mechanism for providing protection
  - ⌘still need TLB lookup or physical tag for writeback
- ⌘Suffer from synonyms and homonyms
  - ⌘requires flushing on context switches
    - ⌘makes context switches expensive
    - ⌘may even be required on kernel→user switch
  - … or guarantee no synonyms and homonyms
- Used on MC68040, i860, ARM7/ARM9/StrongARM/Xscale
- Used for I-caches on several other architectures (Alpha, Pentium 4)
- Not used on recent architectures

Historically used with shallow hierarchies to support bigger L1

Summary: ASID-Tagged VV Caches

- Add ASID as part of tag
- On access, compare with CPU’s ASID register
- ☑Removes homonyms
  - ☑potentially better context-switching performance
  - ⌘ASID recycling still needs flush
- ⌘Doesn’t solve synonym problem (but that’s less severe)
- ⌘Doesn’t solve write-back problem
- Not used on recent architectures

Summary: VP Caches

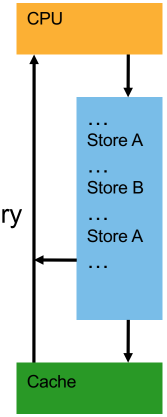
- Medium speed
  - ☑lookup in parallel with address translation
  - ⌘tag comparison after address translation
- ☑No homonym problem
- ⌘Potential synonym problem
- ⌘Bigger tags (cannot leave off set-number bits)
  - ⌘increases area, latency, power consumption
- Used on most contemporary architectures for L1 cache
  - but mostly single-colour (pseudo-PP) or with HW alias prevention (Arm)

Summary: PP Caches

- ⌘Slowest
  - ⌘requires result of address translation before lookup starts
- ☑No synonym problem
- ☑No homonym problem
- ☑Easy to manage
- ☑Cache can use *bus snooping* for DMA/multicore coherency
- ☑Obvious choice for L2–LLC where speed matters less

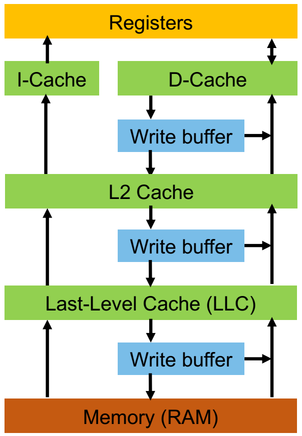
Write Buffer

- Store operations can take a long time to complete
  - eg if a cache line must be read or allocated
- Can avoid stalling the CPU by buffering writes
- *Write buffer* is a FIFO *queue of incomplete stores*
  - Also called *store buffer* or *write-behind buffer*
  - May exist at any cache level, or between cache and memory
- Can fetch intermediate values out of buffer
  - to service read of a value that is still in write buffer
  - avoids unnecessary stalls of load operations
- Implies that memory contents are temporarily stale
  - on a multiprocessor, CPUs see different order of writes!
  - "*weak memory ordering*", to be revisited in SMP context

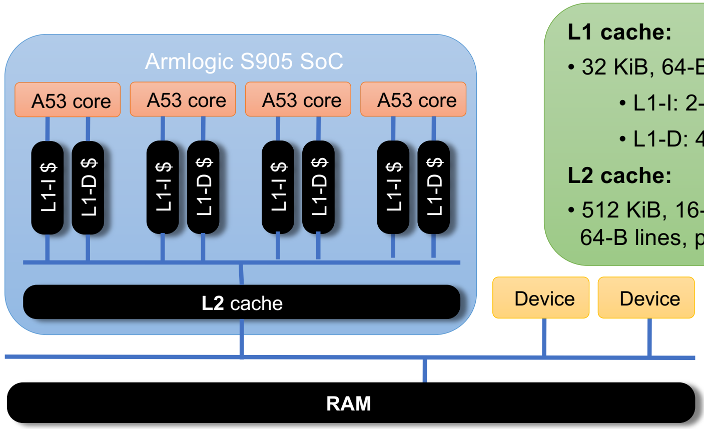


Cache Hierarchy

- Hierarchy of caches to balance memory accesses:
  - small, fast, virtually-indexed L1
  - large, slow, physically indexed L2–LLC
- Each level reduces and clusters traffic
- L1 split into I- and D-caches
  - "Harvard architecture"
  - requirement of pipelining
- Other levels unified
- Chip multiprocessors:
  - Usually LLC shared chip-wide
  - L2 private (Intel) or clustered (AMD)



ODROID-C2 (Cortex A53) System Architecture

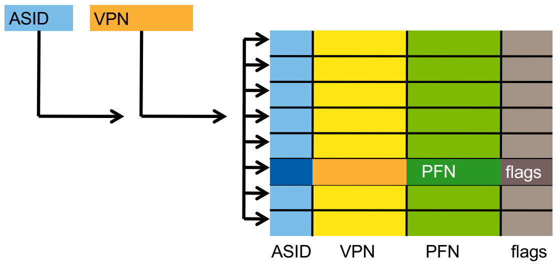


L1 cache:
- 32 KiB, 64-B lines
  - L1-I: 2-way, virtually addr.
  - L1-D: 4-way, "physically" addr.

L2 cache:
- 512 KiB, 16-way
  64-B lines, physical

Translation Lookaside Buffer (TLB)

- TLB is a (VV) cache for page-table entries
- TLB can be
  - hardware loaded, transparent to OS
  - software loaded, maintained by OS
- TLB can be:
  - split: I- and D-TLBs
  - unified



TLB Size (I-TLB+D-TLB)

Not much growth in 40 years!

Architecture	Size (I+D)	Assoc	Page Size	Coverage
VAX-11	64–256	2	0.5 KiB	32–128 KiB
ix86	32i + 64d	4	4 KiB + 4 MiB	128 KiB
MIPS	96–128	full	4 KiB – 16 MiB	384–512 KiB
SPARC	64	full	8 KiB – 4 MiB	512 KiB
Alpha	32–128i + 128d	full	8 KiB – 4 MiB	256 KiB
RS/6000 (PPC)	32i + 128d	2	4 KiB	256 KiB
Power-4 (G5)	1024	4	4 KiB	512 KiB
PA-8000	96i + 96d	full	4 KiB – 64 MiB	384 KiB
Itanium	64i + 96d	full	4 KiB – 4 GiB	384 KiB
ARMv7 (A9)	64–128	1–2	4 KiB – 16 MiB	256–512 KiB
x86 (Skylake)	L1:128i+64d; L2:1536	4	4 KiB + 2/4 MiB	1 MiB

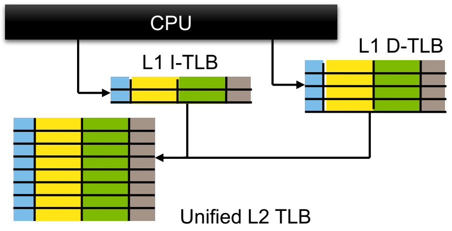
TLB Size

TLB coverage

- Memory sizes are increasing
- Number of TLB entries are roughly constant
- Base page sizes are steady
  - 4 KiB (SPARC, Alpha used 8KiB)
  - OS designers have trouble using superpages effectively
- Consequences:
  - Total amount of RAM mapped by TLB is not changing much
  - Fraction of RAM mapped by TLB is shrinking dramatically!
  - Modern architectures have very low TLB coverage!
- The TLB can become a bottleneck

Multi-Level TLBs

- Multi-level design (like I/D cache)
- Improve size-performance tradeoff



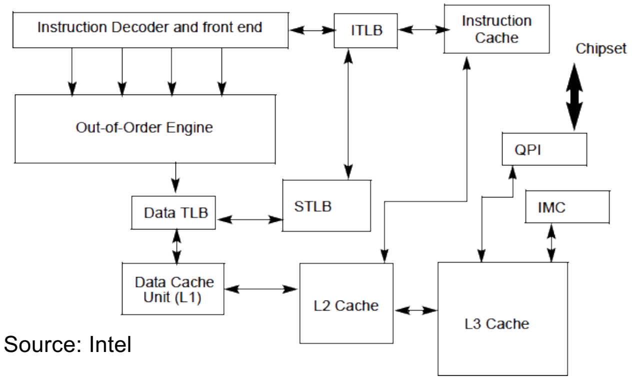
Intel Core i7

L	I/D	Pages	Assoc	Entr
1	I	4 KiB	4-way	64
1	D	4 KiB	4-way	64
1	I	2/4 MiB	fully	7
1	D	2/4 MiB	4-way	32
2	unif	4 KiB	4-way	512

Arm A53

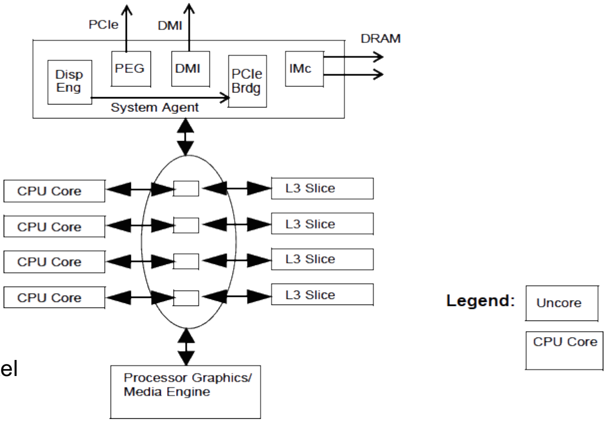
L	I/D	Pages	Assoc	#
1	*I*	*4 KiB–1 GiB?*	*full?*	*10*
1	*D*	*4 KiB–1 GiB?*	*full?*	*10*
2	unif	4 KiB–512 MiB	4-way	512

Intel Core i7 (Haswell) Cache Structure



Source: Intel

Intel Haswell L3 Cache



Source: Intel